## **Riehl-Verity 2.1-2.2 Summary and Questions**

## Brandon Shapiro

These chapters begin the generalization of traditional 1-category theory to arbitrary  $\infty$ -cosmoi, where the 2-category of categories, functors, and natural transformations are replaces by the homotopy 2-category of the cosmos.

## 1. 2.1 : Adjoints

**Definition 1.** An adjunction between  $\infty$ -categories consists of a pair of  $\infty$ -categories A, B, a pair of functors  $u : A \to B$ ,  $f : B \to A$ , and a pair of natural transformations  $\eta : 1_B \Rightarrow uf$  and  $\epsilon : fu \Rightarrow 1_A$  which satisfy the triangle identities in the homotopy 2-category.

Note that this definition is precisely that of an adjunction in the homotopy 2-category, though there is some (purposeful!) vagueness in what is meant by natural transformation in this context (does it exist in the cosmos or its homotopy 2-category).

**Question 2.** What does this look like explicitly in the functor spaces of a cosmos, and how does it simplify possible definitions of an adjunction in that context?

Lemma 3. Adjunctions in a 2-category are preserved by any 2-functor.

Note that this follows directly from the definition of adjunction using unit and counit rather than the alternative definition by isomorphisms of *Hom* sets.

**Example 4.** Take any adjunction of 1-categories and apply the nerve 2-functor to get an adjunction of quasicategories.

**Proposition 5.** Any adjunction in an  $\infty$ -cosmos  $\mathcal{K}$  is preserved by the following 2-functors:

- $Fun(X, -): K \to QCat \text{ for any } \infty\text{-category } X \text{ in } \mathcal{K}$
- $hFun(X, -): K \to Cat \text{ for any } \infty\text{-category } X \text{ in } \mathcal{K}$
- $(-)^U : \mathcal{K} \to \mathcal{K}$  for any simplicial set U
- $(-)^C : \mathcal{K} \to \mathcal{K}$  for any  $\infty$ -category C in  $\mathcal{K}$  if  $\mathcal{K}$  is cartesian closed

**Proposition 6.** For adjunctions  $f : B \rightleftharpoons A : u$  and  $f' : C \rightleftharpoons B : u'$  the composites  $ff' : C \rightleftharpoons A : u'u$  form an adjunction.

**Proposition 7.** Any two left adjoints of  $u : A \to B$  are isomorphic, and any 1-cell isomorphic to a left adjoint of u is also a left adjoint of u.

**Proposition 8.** Any equivalence  $f : A \rightleftharpoons B : g$  can be promoted to an adjoint equivalence by modifying just one of the 2-cells that define it.

Proposition 9. Adjunctions are preserved and reflected by equivalences.

**Lemma 10.** For any  $\infty$ -category A, the composition functor  $A^{\nvDash} \times_A A^{\nvDash} \xrightarrow{\circ} A^{\nvDash}$  admits left and right adjoints which pair an arrow with an identity on the left or right respectively.

Note that all of these results are purely 2-categorical; no additional structure of a cosmos goes into them (other than the cotensoring in the previous lemma).

## 2. Initial and Terminal Elements

We now begin with an example of defining (co)limits in this formalism. Instead of reusing the word "object" to no end, we call maps  $a : 1 \to A$  elements of the  $\infty$ -category A instead.

**Definition 11.** An initial element in an  $\infty$ -category A is a left adjoint to the unique functor  $!: A \to 1$ , and a terminal element is a right adjoint to !.

**Question 12.** Why does this make sense as a definition of initial and terminal objects?

**Lemma 13.** To define an initial element in A it suffices to specify an element  $i : 1 \to A$ and a natural transformation  $\epsilon : i! \Rightarrow 1_A$  such that the component  $\epsilon i : i \Rightarrow i$  is the identity in hA.

**Lemma 14.** An element  $i : 1 \to A$  is initial if and only if for all  $f : X \to A$  there exists a unique 2-cell  $i! \Rightarrow f$ .

This shows an initial element of A to be in fact representably initial, and in particular specializes to show it is also initial among elements in the homotopy category hA.

Lemma 15. Equivalences preserve initial elements up to isomorphism.

*Proof.* An equivalence can be promoted to an adjunction, and adjunctions compose, so if we have an adjunction  $i: 1 \rightleftharpoons A :!$  and an equivalence  $f: A \rightleftharpoons A' : g$ , we get an adjunction  $fi: 1 \rightleftharpoons A' :!g$ .  $\Box$