

18.510: INTRODUCTION TO MATHEMATICAL LOGIC  
AND SET THEORY, FALL 08

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1. PREDICATE CALCULUS, FIRST-ORDER LOGIC

**Semantics.**

*Homomorphism and Isomorphism.* Let  $\mathcal{A}$  and  $\mathcal{B}$  be  $\mathcal{S}$ -structures. A map  $h: A \rightarrow B$  is called an *homomorphism* from  $\mathcal{A}$  to  $\mathcal{B}$  if

- For  $n$ -ary  $R \in \mathcal{S}$  and  $a_1, \dots, a_n \in A$ ,

$$R^{\mathcal{A}}(a_1, \dots, a_n) \text{ iff } R^{\mathcal{B}}(h(a_1), \dots, h(a_n)).$$

- For  $n$ -ary  $f \in \mathcal{S}$  and  $a_1, \dots, a_n \in A$ ,

$$h(f^{\mathcal{A}}(a_1, \dots, a_n)) = f^{\mathcal{B}}(h(a_1), \dots, h(a_n)).$$

- For  $c \in \mathcal{S}$ ,

$$h(c^{\mathcal{A}}) = c^{\mathcal{B}}.$$

If  $h$  is also one-to-one and onto we call it an *isomorphism*, and write  $h: \mathcal{A} \cong \mathcal{B}$ .  $\mathcal{A}$  and  $\mathcal{B}$  are said to be *isomorphic* if there is an isomorphism  $h: \mathcal{A} \cong \mathcal{B}$ .

For example, the  $\mathcal{S}_{\text{gr}}$ -structure  $(\mathbb{N}, +, 0)$  is isomorphic to the  $\mathcal{S}_{\text{gr}}$ -structure  $(G, +^G, 0)$  consisting of the even natural numbers with ordinary addition as  $+^G$ . The map  $\pi: \mathbb{N} \rightarrow G$  mapping  $\pi(n) = 2n$  is an isomorphism of  $(\mathbb{N}, +, 0)$  onto  $(G, +^G, 0)$ .

**1.1. Lemma** (Isomorphism Lemma). *If  $\mathcal{A}$  and  $\mathcal{B}$  are isomorphic  $\mathcal{S}$ -structures then for all  $\mathcal{S}$ -sentences  $\phi$*

$$\mathcal{A} \models \phi \text{ iff } \mathcal{B} \models \phi.$$

You will prove this lemma in PS4.

**1.2. Remark.** However, two  $\mathcal{S}$ -structures might satisfy the same  $\mathcal{S}$ -sentences but not be isomorphic. We will see an example later in the course.

Let  $\mathcal{A}$  and  $\mathcal{B}$  be  $\mathcal{S}$ -structures.  $\mathcal{A}$  is called a *substructure* of  $\mathcal{B}$ , written  $\mathcal{A} \subset \mathcal{B}$ , if  $A \subset B$  and the inclusion map  $i: A \hookrightarrow B$  ( $i(a) = a$ ) is an homomorphism.

For example,  $(\mathbb{Z}, +, 0)$  is a substructure of  $(\mathbb{Q}, +, 0)$ , and  $(\mathbb{N}, +, 0)$  is a substructure of  $(\mathbb{Z}, +, 0)$  (although  $(\mathbb{N}, +, 0)$  is not a subgroup of  $(\mathbb{Z}, +, 0)$ ). Notice that the sentence

$$(\forall x)(\exists y)y + y = x$$

holds in  $(\mathbb{Q}, +, 0)$ , but not in  $(\mathbb{N}, +, 0)$ . So sentences might no longer hold when passing to substructures.

**1.3. Lemma.** *Let  $\mathcal{A}$  and  $\mathcal{B}$  be  $\mathcal{S}$  structure with  $\mathcal{A} \subset \mathcal{B}$  and let  $p: \{v_n\} \rightarrow A$  be an assignment in  $\mathcal{A}$ . Then for every quantifier-free  $\mathcal{S}$ -formula  $\phi$ :*

$$\mathcal{A} \models \phi(p) \text{ iff } \mathcal{B} \models \phi(p).$$

You will prove this lemma in PS4.