The Rotor-Router Shape is Spherical

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In the two-dimensional rotor-router walk (defined by Jim Propp and presented beautifully in [4]), the first time a particle leaves a site x it departs east; the next time this or another particle leaves x it departs south; the next departure is west, then north, then east again, etc. More generally, in any dimension $d \ge 1$, for each site $x \in \mathbb{Z}^d$ fix a cyclic ordering of its 2d neighbors, and require successive departures from x to follow this ordering. In rotor-router aggregation, we start with n particles at the origin; each particle in turn performs rotor-router walk until it reaches an unoccupied site. Let A_n denote the shape obtained from rotor-router aggregation of n particles in \mathbb{Z}^d ; for example, in \mathbb{Z}^2 with the ordering of directions as above, the sequence will be begin $A_1 = \{\mathbf{0}\}$, $A_2 = \{\mathbf{0}, (1, 0)\}, A_3 = \{\mathbf{0}, (1, 0), (0, -1)\}$, etc. As noted in [4], simulations in two dimensions indicated that A_n is close to a ball, but there was no theorem explaining this phenomenon.

Order the points in the lattice \mathbb{Z}^d according to increasing distance from the origin, and let B_n consist of the first n points in this ordering; we call B_n the *lattice ball* of cardinality n. In this letter we outline a proof that for all d, the rotor-router shape A_n in \mathbb{Z}^d is indeed close to a ball, in the sense that

the number of points in the symmetric difference $A_n \Delta B_n$ is o(n). (1)

See [6] for a complete proof, and error bounds. Let $B \subset \mathbb{R}^d$ denote a ball of unit volume centered at the origin, and let $A_n^* \subset \mathbb{R}^d$ be the union of unit cubes centered at the points of A_n ; then (1) means that the volume of the symmetric difference $n^{-1/d}A_n^*\Delta B$ tends to zero as $n \to \infty$. A novel feature of our argument is the use of random walk and Brownian motion to analyze a deterministic cellular automaton.

A stochastic analogue of the rotor-router walk, called *internal diffusion limited aggregation* (IDLA) was introduced earlier by Diaconis and Fulton [3]. In IDLA one also starts with n particles at the origin **0**, and each particle in turn walks until it reaches an unoccupied site; however, the particles perform simple random walk instead of rotor-router walks. Lawler, Bramson and Griffeath [5] showed that the asymptotic shape of IDLA is a ball. Our result does not rely on theirs, but we do use a modification of IDLA in our analysis.

Since the lattice ball B_n minimizes the quadratic weight $Q(A) = \sum_{x \in A} ||x||^2$ among all sets $A \subset \mathbb{Z}^d$ of cardinality n, the difference $Q(A_n) - Q(B_n)$ can be



Figure 1: Rotor-router (left) and IDLA shapes of 10,000 particles. Each site is colored according to the direction in which the last particle left it.

seen as a measurement of how far the set A_n is from a ball. We claim that

$$Q(A_n) \stackrel{<}{\scriptstyle\sim} Q(B_n)$$
 (where $a_n \stackrel{<}{\scriptstyle\sim} b_n$ means that $\limsup a_n/b_n \le 1$). (2)

It is easy to prove that this implies (1). To bound $Q(A_n)$, we use a property of the function $||x||^2$: its value at a point x is one less than its average value on the 2*d* neighbors of x. For a set $A \subset \mathbb{Z}^d$ and a point $x \in \mathbb{Z}^d$, let $\mathcal{E}(x, A)$ be the expected time for random walk started at x to reach the complement of A. If $x \notin A$, then $\mathcal{E}(x, A) = 0$, while if $x \in A$, then $\mathcal{E}(x, A)$ is one more than the average value of $\mathcal{E}(y, A)$ over the 2*d* neighbors y of x. This implies that $h(x) = ||x||^2 + \mathcal{E}(x, A)$ is harmonic in A: its value at $x \in A$ equals its average on the neighbors of x.

Consider rotor-router aggregation starting with n particles at $\mathbf{0}$, and recall that A_n is the set of sites occupied by the particles when they have all stopped. Given a configuration of n particles at (not necessarily distinct) locations x_1, \ldots, x_n , define the harmonic weight of the configuration to be

$$W = W(x_1, \dots, x_n) = \sum_{k=1}^n \left(\|x_k\|^2 + \mathcal{E}(x_k, A_n) \right).$$

We track the evolution of W during rotor-router aggregation. Initially, $W = W(\mathbf{0}, \ldots, \mathbf{0}) = n\mathcal{E}(\mathbf{0}, A_n)$. Since every 2d consecutive visits to a site x result in one particle stepping to each of the neighbors of x, by harmonicity, the net change in W resulting from these 2d steps is zero. Thus the final harmonic weight determined by the n particles, $Q(A_n) + \sum_{x \in A_n} \mathcal{E}(x, A_n)$, equals the initial weight $n\mathcal{E}(\mathbf{0}, A_n)$, plus a small error due to the fact that the number of visits to any given site may not be an exact multiple of 2d. It is not hard to bound this error (see [6]) and deduce that $Q(A_n) \approx n\mathcal{E}(\mathbf{0}, A_n) - \sum_{x \in A_n} \mathcal{E}(x, A_n)$, where $a_n \approx b_n$ means that $\lim a_n/b_n = 1$.

The key step in our argument involves the following modified IDLA. Beginning with n particles $\{p_k\}_{k=1}^n$ at the origin, let each particle p_k in turn



Figure 2: Segments of the boundaries of rotor-router (top) and IDLA shapes formed from one million particles. The rotor-router shape has a smoother boundary.

perform simple random walk until it either exits A_n or reaches a site different from those occupied by p_1, \ldots, p_{k-1} . At the random time τ_n when all the n particles have stopped, the particles that did not exit A_n occupy distinct sites in A_n . If we let these particles continue walking, the expected number of steps needed for all of them to exit A_n is at most $\sum_{x \in A_n} \mathcal{E}(x, A_n)$. Thus $n\mathcal{E}(\mathbf{0}, A_n) \leq \mathbb{E}(\tau_n) + \sum_{x \in A_n} \mathcal{E}(x, A_n)$. So far, we have explained why

$$Q(A_n) \approx n \mathcal{E}(\mathbf{0}, A_n) - \sum_{x \in A_n} \mathcal{E}(x, A_n) \le \mathbb{E}(\tau_n) \,. \tag{3}$$

To estimate $\mathbb{E}(\tau_n)$, we want to bound, for each k < n, the expected number of steps made by the particle p_{k+1} in the random process above; for this, we use a general upper bound on expected exit times from k-point sets in \mathbb{Z}^d . In 1982, Aizenman and Simon [1] showed that among all regions in \mathbb{R}^d of a fixed volume, a ball centered at the origin maximizes the expected exit time for standard d-dimensional Brownian motion started at the origin. (Their proof uses the spherical symmetry of the Gaussian transition density and the powerful Brascamp-Lieb-Luttinger [2] rearrangement inequality.) Since random walk paths are well-approximated by Brownian paths, the Brownian motion result from [1] can be used to prove that for any k-point set $A \subset \mathbb{Z}^d$, the expected exit time $\mathcal{E}(\mathbf{0}, A)$ for random walk is at most $\mathcal{E}(\mathbf{0}, B_k)$ plus a small error term; details may be found in [6]. The number of steps taken by the particle p_{k+1} in our modified IDLA is at most the time for random walk started at $\mathbf{0}$ to exit the set occupied by the stopped particles p_1, \ldots, p_k . It follows that

$$\mathbb{E}(\tau_n) \stackrel{\scriptstyle <}{\scriptstyle \sim} \sum_{k=1}^n \mathcal{E}(\mathbf{0}, B_k) \,. \tag{4}$$

The final step in our argument is to show that $\sum_{k=1}^{n} \mathcal{E}(\mathbf{0}, B_k)$ is approximately equal to $Q(B_n)$. Fix $k \leq n$ and let a single particle p perform random

walk starting at **0** and stopping at the first time t_k that p exits B_k . If S(j) is the location of p after j steps, then the expectation of $||S(j+1)||^2$ given S(j) equals $||S(j)||^2 + 1$. Therefore

$$\mathcal{E}(\mathbf{0}, B_k) = \mathbb{E}(t_k) = \mathbb{E}\Big(\|S(t_k)\|^2\Big).$$
(5)

(Formally, this follows from the Optional Stopping Theorem for Martingales.)

Let v_1, v_2, \ldots be an ordering of \mathbb{Z}^d in increasing distance from the origin, and recall that $B_k = \{v_1, \ldots, v_k\}$. Since all points on the boundary of B_k are about the same distance from the origin, $\mathbb{E}(\|S(t_k)\|^2) \approx \|v_k\|^2$. Summing this over $k \leq n$ and using (5) gives

$$\sum_{k=1}^n \mathcal{E}(\mathbf{0}, B_k) \approx \sum_{k=1}^n \|v_k\|^2 = Q(B_n).$$

Together with (3) and (4), this yields $Q(A_n) \stackrel{<}{\sim} Q(B_n)$, as claimed.

Concluding Remark. As discovered by Jim Propp, simulations in two dimensions indicate that the shape generated by the rotor-router walk is significantly rounder than that of IDLA. One quantitative way of measuring roundness is to compare *inradius* and *outradius*. The inradius of a region A is the minimum distance from the origin to a point not in A; the outradius is the maximum distance from the origin to a point in A. In our simulation up to a million particles, the difference between the inradius and outradius of the IDLA shape rose as high as 15.2. By contrast, the largest deviation between inradius and outradius for the rotor-router shape up to a million particles was just 1.74. Not only is this much rounder than the IDLA shape, it's about as close to a perfect circle as a set of lattice points can get!

Due to error terms incurred along the way, our argument only shows that the rotor-router shape is roughly spherical. It remains a challenge to explain the almost perfectly spherical shapes encountered in simulations.

References

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